

$$C_2 \leq C_0(\mathcal{O}(2)) - c_1 \left(1 + \frac{1 + \log(12\alpha)}{\log(n/\delta)} \right) \cdot \boxed{J}$$

$m \propto S \log(n)$.

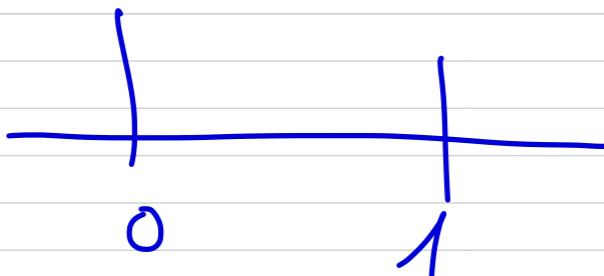
Chapter 10. Approximation Theory

Min-Max Kolmogorov Rate-Distortion Theory

function class $C \subset L^2(\mathcal{S})$, $\mathcal{S} \subset \mathbb{R}^d$

$$\int_{\mathcal{S}} |f(x)|^2 dx < \infty$$

e.g. $[0,1] \rightarrow l$ bits $\rightarrow 2^l$ different values



Set of binary encoders of length ℓ

$$\mathcal{E}^\ell := \{E: C \rightarrow \{0,1\}^\ell\}$$

Set of binary decoders of length ℓ

$$\mathcal{D}^\ell := \{D: \{0,1\}^\ell \rightarrow L^2(\mathbb{R})\}$$

Uniform error ϵ over the function class C

$$\sup_{f \in C} \|D(Ef) - f\|_{L^2(\mathbb{R})} \leq \epsilon.$$

Def. 10.1. Let $d \in \mathbb{N}$, $\Sigma \subset \mathbb{R}^d$, and $C \subset L^2(\mathbb{R})$. Then, for $\epsilon > 0$,
the minimax code length $L(\epsilon, C)$ is

$$L(\epsilon, C) := \min \{l \in \mathbb{N}: \exists (E, D) \in \mathcal{E}^l \times \mathcal{D} : \sup_{f \in C} \|D(Ef) - f\|_{L^2(\mathbb{R})} \leq \epsilon\}.$$

Moreover, the optimal exponent $\eta^*(c)$ is defined as

$$\eta^*(c) = \sup \{ \eta \in \mathbb{R} : L(\varepsilon, c) \in O(\varepsilon^{-1/\eta}), \varepsilon > 0 \}.$$

lower $\eta \Rightarrow$ smaller growth rate!

10.2. Metric entropy, covering, and packing

(\mathcal{X}, ρ)

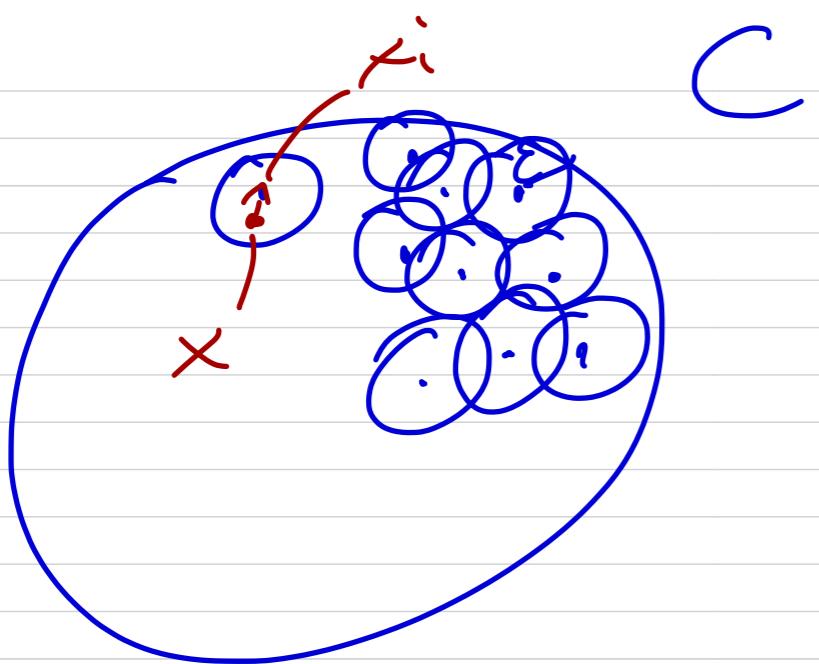
$\rho : \mathcal{X} \times \mathcal{X} \rightarrow \mathbb{R}$

1. $\rho(x, x') \geq 0, \forall x, x' \in \mathcal{X}$, with $=$ if $x = x'$

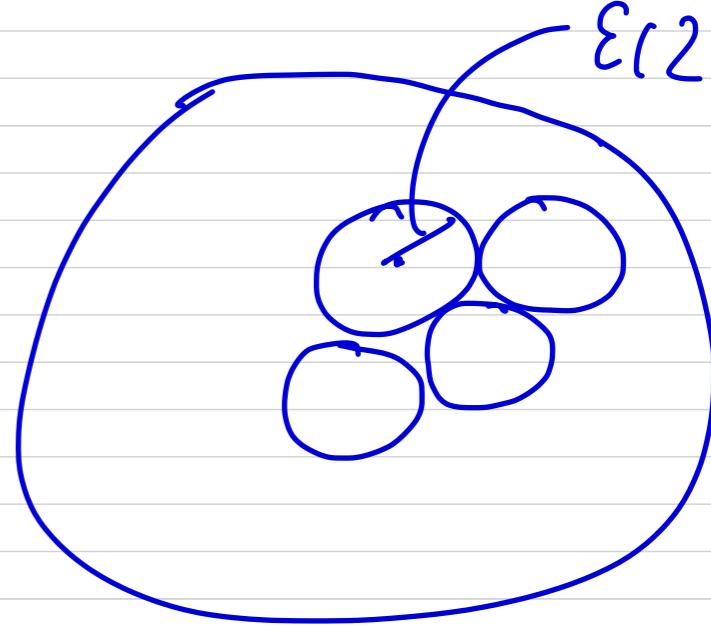
2. $\rho(x, x') = \rho(x', x)$

3. Δ -ineq.: $\rho(x, \tilde{x}) \leq \rho(x, x') + \rho(x', \tilde{x})$, for all x, x', \tilde{x}

$$L^2(C_0, D), \quad \|f - g\|_{L^2(C_0, D)} = \left[\int_0^1 (f(x) - g(x))^2 dx \right]^{1/2}$$



$$\rho(x_i, x_i) \leq \varepsilon$$



covering

Def. 10.2. Let (X, ρ) be a metric space. An ε -covering of $C \subseteq X$ w.r.t. the metric ρ is a set $\{x_1, \dots, x_n\} \subseteq C$ s.t. for each $x \in C$, there exists an $i \in \{1, \dots, n\}$ so that $\rho(x, x_i) \leq \varepsilon$. The ε -covering number $N(\varepsilon; C, \rho)$ is the cardinality of the smallest ε -covering.

$$C \subseteq \bigcup_{i=1}^N B(x_i, \varepsilon)$$

encoder: $x \in C$, map x to the closest (in ℓ) ball center x_i ,
produce a binary representation of x by assigning it
the binary representation of x_i

$N(\varepsilon; C, \ell)$ is the no. of ball centers

metric entropy $\rightarrow \log_2 N(\varepsilon; C, \ell)$ bits to label (in binary form)
The ball centers

decoder: takes bit string of length $\log_2 N(\varepsilon; C, \ell)$ and
maps it to the cent. x_i

$$D(E(x)) = x_i$$

Ex. $\log_2(N; \varepsilon, \ell) \leq \log_2\left(\frac{1}{\varepsilon} + 1\right) \approx \log_2(\varepsilon^{-1})$

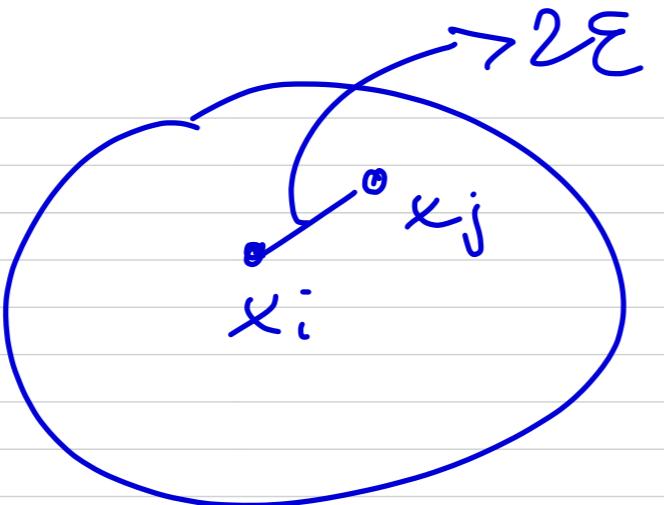
Def. 10.3. Let (X, d) be a metric space. An ε -packing for a compact set $C \subset X$ w.r.t. the metric d is a set $\{x_1, \dots, x_N\} \subset C$ s.t. $d(x_i, x_j) > \varepsilon$, for all $i \neq j$. The ε -packing number $M(\varepsilon; C, d)$ is the cardinality of the largest ε -packing.

Lemma 10.4. Let (X, d) be a metric space and C a compact set in X . For all $\varepsilon > 0$, the packing and the covering numbers are related according to

$$M(2\varepsilon; C, d) \leq N(\varepsilon; C, d) \leq M(\varepsilon; C, d)$$

Proof. Choose minimal ε -covering and a maximal 2ε -packing of C .

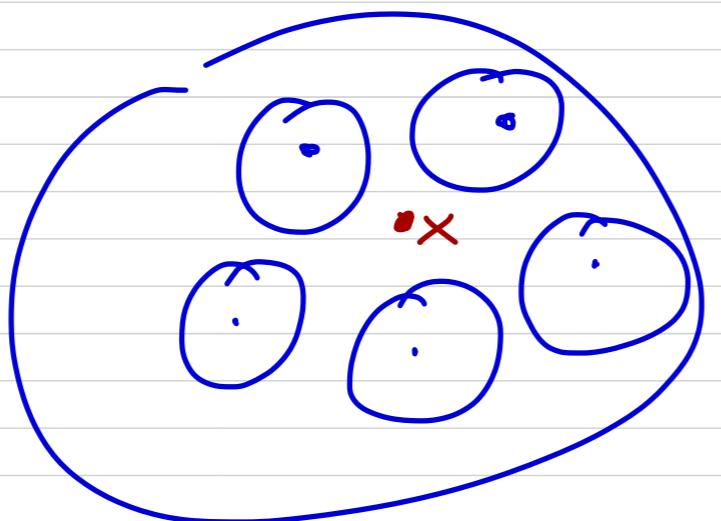
1. no two centers of the 2ε -packing can lie in the same ball of the ε -covering.



$$M(2\epsilon; c_1 \epsilon) \leq N(\epsilon; c_1 \epsilon)$$

2. $N(\epsilon; c_1 \epsilon) \leq M(\epsilon; c_1 \epsilon)$: given an ϵ -packing $M(\epsilon; c_1 \epsilon)$

For given $x \in C$, we have the center of at least one of the balls of the ϵ -packing within dist. $\leq \epsilon$.



\Rightarrow ϵ -packing is also an ϵ -covering. \square

$$M(\epsilon; c_1 \epsilon) \asymp \log_2\left(\frac{1}{\epsilon}\right)$$

$$M(2\epsilon) \leq N(\epsilon) \leq M(\epsilon) \Rightarrow N \asymp \log_2\left(\frac{1}{\epsilon}\right)$$

Lemma 10.5. Consider $\|\cdot\|, \|\cdot\|'$ on \mathbb{R}^d , and let $B \otimes B'$
be their conv. unit balls, i.e.,

$$B = \{x \in \mathbb{R}^d \mid \|x\| \leq 1\}$$

$$B' = \{x \in \mathbb{R}^d \mid \|x\|' \leq 1\}$$

Then, the ε -covering number of B in the $\|\cdot\|'$ -norm
satisfies

$$\left(\frac{1}{\varepsilon}\right)^d \frac{\text{vol}(B)}{\text{vol}(B')} \leq N(\varepsilon; B, \|\cdot\|') \leq \frac{\text{vol}\left(\frac{2}{\varepsilon}B + B'\right)}{\text{vol}(B')}$$

Proof. Let $\{x_1, \dots, x_{N(\varepsilon; B, \|\cdot\|')}\}$ be an ε -covering
of B in the $\|\cdot\|'$ -norm. Then,

$$N(\varepsilon; B, \|\cdot\|')$$

$$B \subseteq \bigcup_{j=1}^N \{x_i + \varepsilon B'\} \quad | \text{ vol}(\cdot)$$

$$\text{vol}(\mathcal{B}) \leq N(\varepsilon; \mathcal{B}, \|\cdot\|') \underbrace{\text{vol}(\varepsilon \mathcal{B}^1)}_{\varepsilon^d \text{ vol}(\mathcal{B}^1)}$$

$$\Rightarrow N(\varepsilon; \mathcal{B}, \|\cdot\|') \geq \frac{\text{vol}(\mathcal{B})}{\text{vol}(\mathcal{B}^1)} \left(\frac{1}{\varepsilon}\right)^d \checkmark$$

maximal ε -packing $\{x_1, \dots, x_{N(\varepsilon; \mathcal{B}, \|\cdot\|')}\}$ of \mathcal{B} in the $\|\cdot\|'$ -norm. The balls $\{x_j + \frac{\varepsilon}{2} \mathcal{B}^1, j=1, \dots, N(\varepsilon; \mathcal{B}, \|\cdot\|')\}$ are disjoint and contained within $\mathcal{B} + \frac{\varepsilon}{2} \mathcal{B}^1$.

Taking volumes, we get

$$\sum_{j=1}^{N(\varepsilon; \mathcal{B}, \|\cdot\|')} \text{vol}(x_j + \frac{\varepsilon}{2} \mathcal{B}^1) \leq \text{vol}(\mathcal{B} + \frac{\varepsilon}{2} \mathcal{B}^1)$$

$$= N(\varepsilon; \mathcal{B}, \|\cdot\|') \underbrace{\text{vol}(\frac{\varepsilon}{2} \mathcal{B}^1)}_{1} \leq \text{vol}(\frac{\varepsilon}{2} (\frac{2}{\varepsilon} \mathcal{B} + \mathcal{B}^1))$$

$$\left(\frac{\varepsilon}{2}\right)^d \text{vol}(\mathcal{D})$$

~~$\left(\frac{\varepsilon}{2}\right)^d \text{vol}$~~
 $\left(\frac{2}{\varepsilon}\mathcal{B} + \mathcal{B}^1\right)$

$$H(\varepsilon; \mathcal{B}, \|\cdot\|) \geq \alpha(\varepsilon; \mathcal{B}, \|\cdot\|)$$

$$N(\varepsilon; \mathcal{D}, \|\cdot\|) \leq \frac{\text{vol}\left(\frac{2}{\varepsilon}\mathcal{B} + \mathcal{B}^1\right)}{\text{vol}(\mathcal{B}^1)} \cdot D$$

metric entropy of unit ball in its own norm

$$\begin{aligned} \mathcal{B}^1 = \mathcal{B} : \text{vol}\left(\frac{2}{\varepsilon}\mathcal{B} + \mathcal{B}^1\right) &= \text{vol}\left(\left(\frac{2}{\varepsilon} + 1\right)\mathcal{B}\right) \\ &= \left(\frac{2}{\varepsilon} + 1\right)^d \text{vol}(\mathcal{B}) \end{aligned}$$

$$\left(\frac{1}{\varepsilon}\right)^d \leq N(\varepsilon; \mathcal{B}, \|\cdot\|) \leq \left(\frac{2}{\varepsilon} + 1\right)^d \frac{\text{vol}(\mathcal{B})}{\text{vol}(\mathcal{B})}$$

$$\mathcal{B}: \quad N(\varepsilon; \mathcal{B}, \|\cdot\|) \asymp \varepsilon^{-d}$$

$$[0,1]^d: \quad N(\varepsilon; \mathcal{B}, \|\cdot\|) \asymp d \log(1/\varepsilon)$$

Lipschitz functions: $\mathcal{F}_L := \{g: [0,1] \rightarrow \mathbb{R} \mid g(0)=0, |g(x)-g(x')| \leq L|x-x'|, \forall x, x' \in [0,1]\}.$

$$\log_2(N(\varepsilon; \mathcal{F}_L, \|\cdot\|_\infty)) \asymp L/\varepsilon.$$

$$\mathcal{F}_L([0,1]^d): \quad \log_2(N(\varepsilon; \mathcal{F}_L, \|\cdot\|_\infty)) \asymp (L/\varepsilon)^d$$

$$L(\varepsilon; C) \in O(\varepsilon^{-1/n}) \leq C^d \varepsilon^{-1/n}$$

$$\text{General scaling behavior: } \log_2(N(\varepsilon; C, \|\cdot\|)) \asymp \varepsilon^{-1/n} \log\left(\frac{1}{\varepsilon}\right)$$

ϵ
or
Some
functions
scaling
slower
than
 $\epsilon^{-1/n}$

$$\epsilon^{-1/n} \log^{\beta}(\epsilon^{-1}) \in O(\epsilon^{-1/(n+\beta)})$$

10.3. Approximation with Representation Systems

\mathcal{H} ... Hilbert space, $\langle \cdot, \cdot \rangle$, $\|\cdot\|$

$\{\epsilon_k\}_{k=1}^\infty$ an ONB for \mathcal{H}

1. Linear approximation

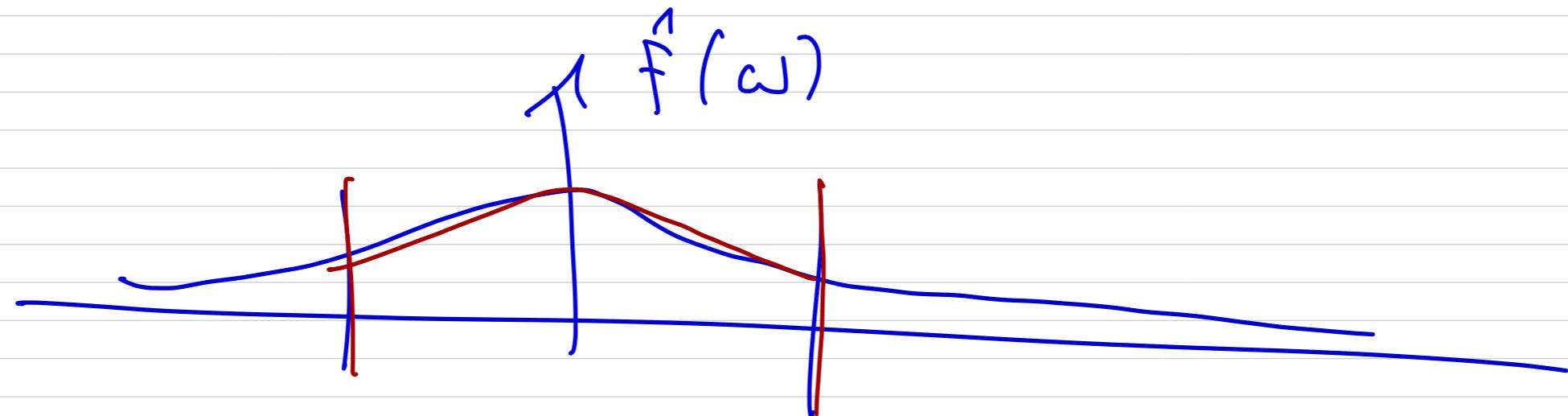
$$\mathcal{H}_M := \text{span} \{\epsilon_k : 1 \leq k \leq M\}$$

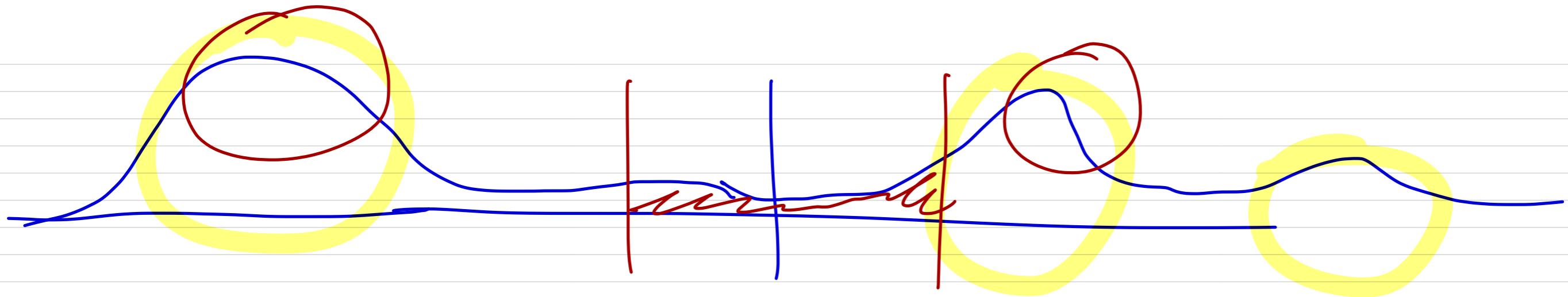
$\bar{E}_M(f)_{\mathcal{H}} := \inf_{g \in \mathcal{H}_M} \|f - g\|_2$.

$$g_1 \in \mathcal{H}_M, g_2 \in \mathcal{H}_M : g_1 = \sum_{g=1}^M c_g^{(1)} g_g$$

$$g_2 = \sum_{g=1}^M c_g^{(2)} g_g$$

$$\alpha g_1 + \beta g_2 = \sum_{g=1}^M (\underbrace{\alpha c_g^{(1)} + \beta c_g^{(2)}}_{n_g}) g_g \in \mathcal{H}_M$$





2. Nonlinear approximation

Nonlinear approximation

replace \mathcal{J}_M by \mathcal{I}_M consisting of all g est that can be expressed as

$$g = \sum_{\Delta \in \Delta} c_\Delta e_\Delta$$

with Δ CN s.t. $|\Delta| \leq M$.