

Thm. 10.8. $p^{*\text{eff}}(C, D) \leq p^*(C)$

Proof. effective best n -term approx. of $f \in C$

$$\|f - \sum_{i \in I_m} c_i q_i\| \leq Cn^{-r}$$

$$n < p^{*\text{eff}}(C, D)$$

$$|I_m| \leq \pi(n)$$

1. indices of participating dictionary elements

$$\log_2 \pi(n) \leq \log_2 n^P = P \log_2 n$$

$$M \log_2 M = \underbrace{Cn \log_2 M}_{= O(M \log_2 M)}$$

2. quantization of the c_i

$$\sum_{i \in I_m} c_i q_i = \sum_{i \in \tilde{I}_m} \hat{c}_i \tilde{q}_i$$

$$\tilde{I}_m \leq M$$

$$e = f - \sum_{i \in \tilde{I}_m} \hat{c}_i \tilde{q}_i = f - \sum_{i \in I_m} c_i q_i$$

$$\left\| \sum_{i \in \tilde{I}_m} \hat{c}_i \tilde{q}_i \right\| = \|f - e\| \leq \|f\| + \|e\|$$

$$\underbrace{\sqrt{\sum_{i \in \tilde{I}_m} |\hat{c}_i|^2}}_{\leq \underbrace{\sup_{f \in C} \|f\| + Cn^{-r}}_{\leq D}} = D + Cn^{-r} < \infty$$

Quantization levels : $\tilde{q}_i^{(n+1/2)} \in \mathbb{Z}$

$$\text{no. of bits needed to encode one } \tilde{q}_i^{(n+1/2)} = \log_2 \frac{C'}{\tilde{q}_i^{(n+1/2)}}$$

$$= \omega_2 C' n^{(n+1)/2}$$

$$c_i = \alpha(\tilde{c}_i) \quad (\lfloor \rfloor, \lceil \rceil)$$

$$= C'' \omega_2 M$$

$$\# \text{bits needed to encode all } \tilde{c}_i = C'' M \log_2 M$$

$$\| f - \sum_{i \in \tilde{\Sigma}_M} \tilde{c}_i \hat{q}_i \| = \| f - \sum_{i \in \tilde{\Sigma}_M} \tilde{c}_i \hat{q}_i + \sum_{i \in \tilde{\Sigma}_M} \tilde{c}_i \hat{q}_i - \sum_{i \in \tilde{\Sigma}_M} \tilde{c}_i \hat{q}_i \|$$

$$\leq CM^{-n} \leq \| f - \sum_{i \in \tilde{\Sigma}_M} \tilde{c}_i \hat{q}_i \| + \left\| \sum_{i \in \tilde{\Sigma}_M} (\tilde{c}_i - \hat{c}_i) \hat{q}_i \right\|$$

$$= \| e \| \leq CM^{-n}$$

$$= \left(\sum_{i \in \tilde{\Sigma}_M} | \tilde{c}_i - \hat{c}_i |^2 \right)^{1/2}$$

$$\leq CM^{-n}$$

$$\sum_{i \in \tilde{\Sigma}_M} | \tilde{c}_i - \hat{c}_i |^2 \leq \sum_{i \in \tilde{\Sigma}_M} M^{-2n-1} \leq M \cdot M^{-2n-1}$$

$$= M^{-2n}$$

$$|\tilde{\Sigma}_M| \leq |\Sigma_M| \leq M$$

$$\| f - \sum_{i \in \tilde{\Sigma}_M} \tilde{c}_i \hat{q}_i \| \leq \underbrace{CM^{-n}}_{\epsilon}$$

$$\# \text{bits} = \Theta(M \log_2 M)$$

1. recall that all arguments we used are valid for $\gamma < \gamma^{*\text{eff}(C,D)}$

$$2. \# \text{bits} = \Theta(M \log_2 M) = \Theta(\epsilon^{-1/n} \log_2 \epsilon^{-1/n})$$

$$= \Theta(\epsilon^{-1/(\gamma^{*\text{eff}(C,D)} - \delta)})$$

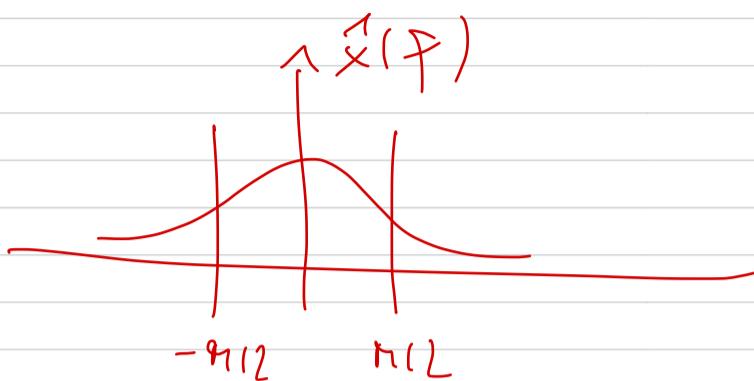
$$\gamma^{*\text{eff}(C,D)} - \delta \leq \gamma^*(C).$$

Def. 10.9.

$$\gamma^{*\text{eff}(C,D)} = \gamma^*(C)$$

C is optimally representable in D

Chapter 11



11.1.1. Linear Approximation

Let $\mathcal{B} = \{g_j\}_{j \in \mathbb{N}_0}$ be an ONS for H . Every $x \in H$ can be written as

$$x = \sum_{j=0}^{\infty} \langle x, g_j \rangle g_j$$

approximate x by the first M terms, we get

$$x_M = \sum_{j=0}^{M-1} \langle x, g_j \rangle g_j$$

$$x - x_M = \sum_{j=M}^{\infty} \langle x, g_j \rangle g_j$$

$$\begin{aligned} E[\epsilon] &= \|x - x_M\|^2 = \left\| \sum_{j=0}^{\infty} \langle x, g_j \rangle g_j - \sum_{j=0}^{M-1} \langle x, g_j \rangle g_j \right\|^2 \\ &= \left\| \sum_{j=M}^{\infty} \langle x, g_j \rangle g_j \right\|^2 = \sum_{j=M}^{\infty} |\langle x, g_j \rangle|^2 \end{aligned}$$

$$\|x\|^2 < \infty$$

$$\sum_{j=0}^{\infty} |\langle x, g_j \rangle|^2 < \infty$$

$\lim_{n \rightarrow \infty} E_e(n) = 0$? how fast

$$j^{2s} j^{-2} = j^{1+2-s} = j^{-1+s}$$

Thm. 11.1. Let $s > 1/2$ and $\sum_{j=0}^{\infty} j^{2s} |\langle x, g_j \rangle|^2 < \infty$. Then there exists constants $A, B > 0$ such that

$$j^{-s-1/2-\varepsilon} \Rightarrow j^{2s} \cdot -2s-1-\varepsilon$$

$$A \sum_{j=0}^{\infty} j^{2s} |\langle x, g_j \rangle|^2 \leq \sum_{n=0}^{\infty} n^{2s-1} E_e(n) \leq B \sum_{j=0}^{\infty} j^{2s} |\langle x, g_j \rangle|^2$$

and hence $E_e(n) = o(n^{-2s})$, i.e., $\lim_{n \rightarrow \infty} E_e(n) n^{2s} = 0$.

$$E_e(n) \sim n^{-2s}$$

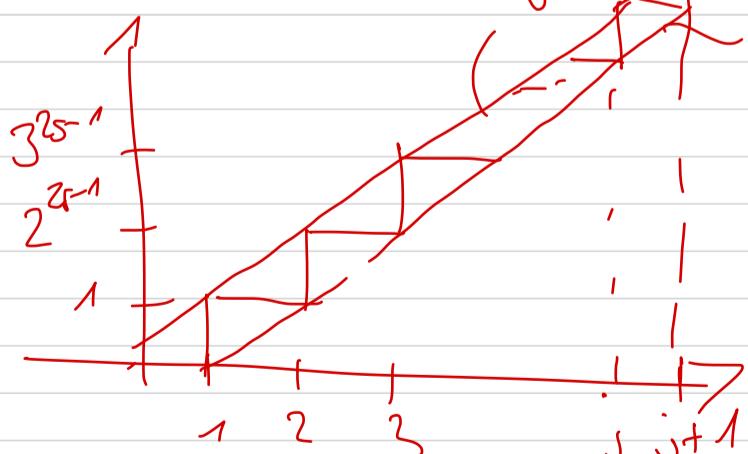
Proof.

$$\sum_{n=0}^{\infty} n^{2s-1} E_e(n) = \sum_{n=0}^{\infty} n^{2s-1} \sum_{j=n}^{\infty} |\langle x, g_j \rangle|^2$$

$$= 1 \cdot \sum_{j=1}^{\infty} |\langle x, g_j \rangle|^2 + 2^{2s-1} \sum_{j=2}^{\infty} |\langle x, g_j \rangle|^2 + \dots$$

$$= \sum_{j=0}^{\infty} |\langle x, g_j \rangle|^2 \sum_{n=0}^j n^{2s-1}$$

$$\underbrace{\int_0^j y^{2s-1} dy}_{\sim j^{2s}} \leq \sum_{n=0}^j n^{2s-1} \leq \underbrace{\int_0^{j+1} y^{2s-1} dy}_{\leq B j^{2s}}$$



Verify that $E_e(n) = o(n^{-2s})$

$$E_e(n) \sum_{j=\lfloor n/2 \rfloor}^{n-1} j^{2s-1} \leq \sum_{j=\lfloor n/2 \rfloor}^{n-1} j^{2s-1} E_e(j) \leq \sum_{j=\lfloor n/2 \rfloor}^{\infty} j^{2s-1} E_e(j)$$

$$\lim_{n \rightarrow \infty} \sum_{j=\lfloor M/2 \rfloor}^{\infty} j^{2s-1} \varepsilon_e[j] = 0$$

$$\sum_{j=\lfloor M/2 \rfloor}^{\infty} j^{2s-1} \geq CM^{2s}$$

$$\lim_{n \rightarrow \infty} (M^{2s} \varepsilon_e[M]) = 0 \Rightarrow \varepsilon_e[M] = o(M^{-2s}) \text{ . q.e.d.}$$

$$W_{\mathcal{B},s} = \left\{ x \in \mathcal{H} : \sum_{j=0}^{\infty} j^{2s} |\langle x, g_j \rangle|^2 < \infty \right\}.$$

ONB

$W_{\mathcal{B},s}$ for $\mathcal{B} = \text{Fourier or wavelet}$ is a Sobolev space.

11.1.2 Regularity & Decay

Theorem 11.2. If $x \in L_1$, then \hat{x} is uniformly continuous and satisfies

$$|\hat{x}(f)| \leq \int_{-\infty}^{\infty} |x(t)| dt < \infty, \quad f \in \mathbb{R}$$

and

$$\lim_{|f| \rightarrow \infty} \hat{x}(f) = 0.$$

$x \in L_1 \Rightarrow x$ is unif. cont. & bounded and satisfies $\lim_{|t| \rightarrow \infty} x(t) = 0$.

$$x'(t) \rightsquigarrow (2\pi i f) \hat{x}(f)$$

$$\hat{x}(f) = \int x(t) e^{-i 2\pi f t} dt$$

Prop. 11.3. A function x is bounded and p times continuously differentiable with bounded derivative if

$$\int_{-\infty}^{\infty} |\hat{x}(f)| / (1+|f|)^p df < \infty$$

Proof. $x^{(q)}(f) \in C_{\text{lf}, f}^q \subset \mathcal{X}(f)$

$$|x^{(q)}(f)| \leq C \int_{-\infty}^{\infty} |x'(f)| |f|^q \leq C \int_{-\infty}^{\infty} |\mathcal{Z}(f)| (1+|f|)^q df$$

$$\Rightarrow \int_{-\infty}^{\infty} |\mathcal{Z}(f)| (1+|f|)^q df < \infty \text{, for every } q \leq p.$$

$$|\mathcal{Z}'(f)| \leq \frac{C'}{(1+|f|)^{p+1+\varepsilon}}$$

III.1.3. Non-linear Approximation

B. - ONB

$$x = \sum_{j=0}^{\infty} \langle x, g_j \rangle g_j \quad | \quad x_n = \sum_{j \in \mathbb{N}_n} \langle x, g_j \rangle g_j$$

$$\mathcal{E}[n] = \|x - x_n\|^2 = \left\| \sum_{j=0}^{\infty} \langle x, g_j \rangle g_j - \sum_{j \in \mathbb{N}_n} \langle x, g_j \rangle g_j \right\|^2$$

$$= \left\| \sum_{j \notin \mathbb{N}_n} \langle x, g_j \rangle g_j \right\|^2 = \sum_{j \notin \mathbb{N}_n} |\langle x, g_j \rangle|^2$$

$$\|x\|^2 = \sum_j |\langle x, g_j \rangle|^2 = \sum_{j \in \mathbb{N}_n} |\langle x, g_j \rangle|^2 + \sum_{j \notin \mathbb{N}_n} |\langle x, g_j \rangle|^2$$

$$\mathcal{E}[n] = \|x\|^2 - \sum_{j \in \mathbb{N}_n} |\langle x, g_j \rangle|^2$$